

# Propositions and rigidity in Layered DRT

Emar Maier and Bart Geurts

University of Nijmegen

<{ | }@phil.kun.nl>

ESLLI 2003 Workshop "Direct Reference and Specificity"

19th August 2003

## Kripke/Kaplan on rigid terms

- names and indexicals are directly referential/rigid designators
- wide-scope behavior w.r.t. operators
- not synonymous with the description giving their 'descriptive meaning', as shown by Kripke-Kaplan examples (1) and (2):
  - 'Sam'  $\neq$  'the person called 'Sam''
    - (1) a. Sam is called Sam
    - b. The person called Sam is called Sam
  - 'I'  $\neq$  'the speaker'
    - (2) a. I am the speaker
    - b. The speaker is the speaker

## Discourse Representation Theory

- meaning encoded in descriptive representational conditions in a DRS
- definite NPs treated as presuppositions
- 2-stage interpretation ( ; ):
  - syntactic module **Prel** builds preliminary DRS from sentence
  - pragmasemantic resolution algorithm
    - \* merges **Prel**( $\sigma$ ) with previous discourse-/background-DRS

- \* performs anaphora and presupposition **Resolution** (binding new presuppositions to old representations or accommodating them)

## Resolution

1. presuppositions are to be bound (e.g. anaphors)
2. otherwise, accommodate as high as possible (e.g. some definite descriptions, factive complements)
  - presupposition resolution accounts for definite NP's wide-scope behavior
  - how about the alleged special behavior of proper names/indexicals?

## Descriptivism

suggests that names are like other definite NP's

- both definite, thus presupposition triggers
- reduction: 'John' = 'the person called 'John''
- accounts for proper name's normal wide scope behavior,
- but also the (rare) other behaviors that proper names appear to share with definite descriptions. . .
- no familiarity required (global accommodation):
  - (3) My best friend is John
- narrow scope (local and intermediate accommodation):
  - (4) If alphabetical order had been the method of electing the American president, Aaron Aardvark might have been president
- bound variable (presupposition binding/cancellation):
  - (5) If a child is christened 'Bambi', then Disney will sue Bambi's parents
- but no account for Kaplan-Kripke examples!

## Layered DRT

- (conservative) extension of standard DRT

- abstract interface for representing the interaction of different *types* of information
- layers for e.g. presuppositions, implicatures, asserted (Fregean) content, syntactic features, . . .
- all layers get truthconditional evaluation

## Applications

- binding problems (reference marker can be employed at several layers at once )
- denial (denial can be directed at one specific layer )
- rigidity

## Rigid designation in LDRT

- represent proper names and indexicals as descriptive conditions,
- but at separate layer '*k*'
- add 2-dimensional semantics to rigidify only that layer
- extension: at representational level, treat *k* like presupposition layer (to account for bound variable uses of names)

## Syntax

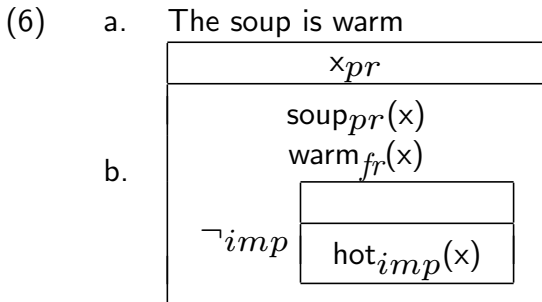
- primitive symbols:
  - a set  $\mathcal{X}$  of reference markers
  - some sets  $\mathcal{P}red^n$  of *n*-place predicates
  - a set  $\Lambda$  of layer labels
- if  $x \in \mathcal{X}$ ,  $L \subseteq \Lambda$ , then  $x_L$  is a labeled reference marker
- if  $x_1, \dots, x_n \in \mathcal{X}$ ,  $L \subseteq \Lambda$ , then  $P_L(x_1, \dots, x_n)$  is a labeled condition
- if  $\varphi$  and  $\psi$  are labeled conditions,  $L \subseteq \Lambda$ , then  $\neg_L \varphi$ ,  $\varphi \vee_L \psi$ , and  $\varphi \Rightarrow_L \psi$  are also labeled conditions

- if  $U$  is a set of labeled reference markers and  $Con$  a set of labeled conditions, then  $\langle U, Con \rangle$  is an LDRS

### Application

- $\Lambda = \{fr, k, imp, pr\}$ 
  - $fr$  = the Frege layer, what is “said”
  - $k$  = Kaplan-Kripke layer for rigid stuff
  - $imp$  = implicature layer
  - $pr$  = (unresolved) presupposition layer

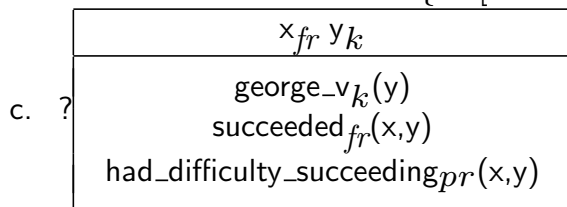
### Examples



### Binding problems

e.g. with presuppositions ( )

- (7) a. ?Someone managed to succeed George V on the throne of England.  
 b.  $\exists x[\text{succeed\_george\_v}(x)]\{\exists x[\text{had\_diff\_succ\_george\_v}(x)]\}$



## Semantics

- $\mathcal{M} = \langle D, W, I \rangle$ 
  - $D$  is a domain of individuals
  - $W$  is a set of possible worlds
  - $I$  interpretation function (from basic predicates to their intensions in  $D^W$ )
- $\|\varphi\|_{L,w}^f$ 

$$\left\{ \begin{array}{ll} = \{g \mid f \subseteq g \wedge \text{Dom}(g) = \\ = \text{Dom}(f) \cup U_L(\varphi) \wedge \\ \wedge \forall \psi \in \text{Con}(\varphi) [\|\psi\|_{L,w}^g = 1]\} & \text{if } \exists g [f \subseteq g \wedge \text{Dom}(g) \\ & = \text{Dom}(f) \cup U_L(\varphi) \wedge \\ & \wedge \forall \psi \in \text{Con}(\varphi) \\ & [\|\psi\|_{L,w}^g \in \{0, 1\}]] \\ \text{undefined} & \text{otherwise} \end{array} \right.$$
- $\|P_K(x_1, \dots, x_n)\|_{L,w}^f$ 

$$\left\{ \begin{array}{ll} = 1 & \text{if } K \cap L = \emptyset \text{ or} \\ & (x_1, \dots, x_n \in \text{Dom}(f) \text{ and} \\ & \langle f(x_1), \dots, f(x_n) \rangle \in I(P)(w)) \\ = 0 & \text{if } K \cap L \neq \emptyset \text{ and} \\ & x_1, \dots, x_n \in \text{Dom}(f) \text{ and} \\ & \langle f(x_1), \dots, f(x_n) \rangle \notin I(P)(w) \\ \text{undefined} & \text{otherwise} \end{array} \right.$$
- $\|\neg_K \psi\|_{L,w}^f$ 

$$\left\{ \begin{array}{ll} = 1 & \text{if } K \cap L = \emptyset \text{ or } \|\psi\|_{K \cap L, w}^f = \emptyset \\ = 0 & \text{if } K \cap L \neq \emptyset \text{ and } \|\psi\|_{K \cap L, w}^f \text{ defined and} \\ & \|\psi\|_{K \cap L, w}^f \neq \emptyset \\ \text{undefined} & \text{otherwise} \end{array} \right.$$
- $\|\psi \vee_K \chi\|_{L,w}^f$ 

$$\left\{ \begin{array}{ll} = 1 & \text{if both } \|\psi\|_{K \cap L, w}^f \text{ and } \|\chi\|_{K \cap L, w}^f \text{ defined,} \\ & \text{and } \|\psi\|_{K \cap L, w}^f \cup \|\chi\|_{K \cap L, w}^f \neq \emptyset \\ = 0 & \text{if } \|\psi\|_{K \cap L, w}^f = \|\chi\|_{K \cap L, w}^f = \emptyset \\ \text{undefined} & \text{otherwise} \end{array} \right.$$

## Layered content

- $\mathcal{C}_L^f(\varphi) = \begin{cases} \{w \in \mathcal{W} \mid \exists g \in \|\varphi\|_{L,w}^f\} & \text{if } \exists w[\|\varphi\|_{L,w}^f \text{ defined}] \\ \text{undefined} & \text{otherwise} \end{cases}$
- $\mathcal{C}_L(\varphi) = \mathcal{C}_L^\emptyset(\varphi)$

## Problems

often undefined (because of layer interdependencies):

- (8) a. The King of France is bald

$x_{pr}$
king_of_france <sub>pr</sub> (x) bald <sub>fr</sub> (x)

- b.  $\mathcal{C}_{pr,fr}(\quad) = \{w \in W \mid \text{there is a bald king of France in } w\}$   
 c.  $\mathcal{C}_{fr}(\varphi)$  undefined

- (9) a. I am speaking

$x_k$
speaker <sub>k</sub> (x) speaking <sub>fr</sub> (x)

- b. The speaker is speaking

$x_{fr}$
speaker <sub>fr</sub> (x) speaking <sub>fr</sub> (x)

- c.  $\mathcal{C}_{k,fr}(\quad) = \mathcal{C}_{k,fr}(\quad) = \mathcal{C}_{fr}(\quad) = W$   
 d.  $\mathcal{C}_{fr}(\quad)$  undefined

to insure definedness, add extra layers as 'background' for interpretation. . .

## The weak proposal

weak one-dimensional content substitute:

$$(10) \quad \mathcal{C}_L^{[K]}(\varphi) = \mathcal{W} - (\mathcal{C}_K(\varphi) - \mathcal{C}_{K \cup L}(\varphi)) = \{w \in \mathcal{W} \mid w \in \mathcal{C}_K(\varphi) \rightarrow w \in \mathcal{C}_{K \cup L}(\varphi)\}$$

handy for denial applications, but:

$$(11) \quad \mathcal{C}_{fr}^{[k]}(\quad) = W$$

## Adding another dimension

- $\mathcal{M} = \langle D, W, C, I \rangle$ 
  - $D$  is a domain of individuals
  - $W$  is a set of worlds
  - $I$  is interpretation function
  - $C$  is a set of contexts, i.e. for all  $c \in C$ :  $I(\text{speaker})(c)$ ,  $I(\text{Mary})(c)$ , . . . are singletons
- $\mathcal{K}_L^c(\varphi) \begin{cases} = \mathcal{C}_L^f(\varphi) & \text{if } \|\varphi\|_{k,c}^\emptyset = \{f\} \\ \text{undefined} & \text{otherwise} \end{cases}$

(12) a. Ashley is called 'Ashley'

$x_k$
ashley <sub>k</sub> (x) ashley <sub>fr</sub> (x)

c. Ashley is Ashley

$x_k \ y_k$
ashley <sub>k</sub> (x) ashley <sub>k</sub> (y) $x =_{fr} y$

d.

$$\mathcal{C}_{k,fr}(\quad) = \mathcal{C}_{k,fr}(\quad) = W, \mathcal{C}_{fr}(\quad), \mathcal{C}_{fr}(\quad) \text{ undefined}$$

$$\mathcal{K}_{fr}^c(\quad) = \mathcal{C}_{fr}^{\{\langle x, \mathbf{ash}_c \rangle\}}(\quad) = \{w \mid \mathbf{ash}_c \text{ is in } w \text{ called 'Ashley'}\}$$

$$\mathcal{K}_{fr}^c(\quad) = \mathcal{C}_{fr}^{\{\langle x, \mathbf{ash}_c \rangle, \langle y, \mathbf{ash}_c \rangle\}}(\quad) = W$$

## Refinements

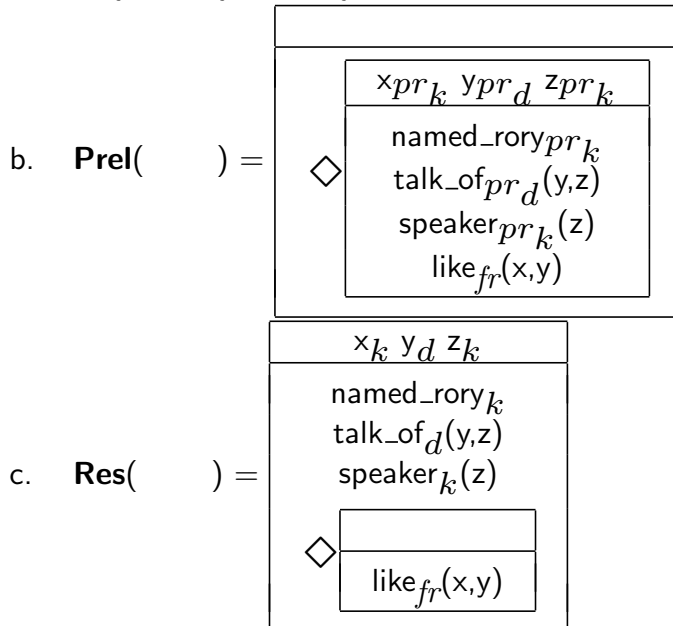
suggests that names are like other definite NP's, with e.g. bound variable uses:

- (13) a. If a child is christened 'Bambi', then Disney will sue Bambi's parents  
 b. Every time we do our Beatles act, Ringo gets drunk afterwards

### Names as presuppositions

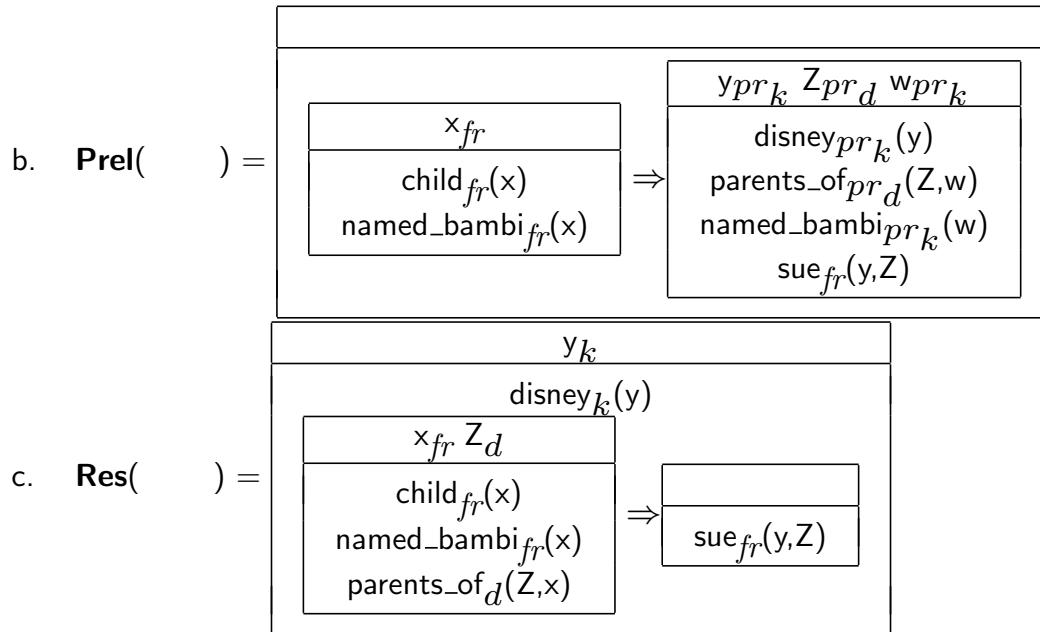
- $\Lambda = \{pr_d, d, pr_k, k, fr, imp\}$ 
  - $pr_d$  = descriptive presupposition layer
  - $d$  = accommodated (descriptive) presupposition
  - $pr_k$  = name/indexical presupposition
  - $k$  = rigidifiable layer
- **Prel** puts names and indexicals in  $pr_k$ , definite descriptions in  $pr_d$
- **Res** tries to bind  $pr_k$  and  $pr_d$  whenever possible, otherwise  $pr_k$  stuff is dropped as high up as possible in  $k$ ,  $pr_d$  gets dropped at  $d$

- (14) a. Maybe Rory liked my talk



- (15) a. If a child is christened 'Bambi', then Disney will sue Bambi's parents





### Further research

- The dynamics of the  $pr_k$  layer: how, why, when does  $pr_k$  stuff become part of new descriptive content (accommodation examples)?
- difference between names and indexicals?

### Appendix I: syntax

The primitive symbols of an LDRT language are:

- a set  $\mathcal{X}$  of reference markers
- some sets  $\mathcal{Pred}^n$  of  $n$ -place predicates
- a set  $\Lambda$  of layer labels

The rest of the syntax is:

- if  $x \in \mathcal{X}$ ,  $L \subseteq \Lambda$ , then  $x_L = \langle x, L \rangle \in \mathcal{X} \times \wp(\Lambda)$  is a labeled reference marker
- if  $P \in \mathcal{Pred}^n$ ,  $L \subseteq \Lambda$ , then  $P_L$  is a labeled predicate
- if  $x, y \in \mathcal{X}$ ,  $L \subseteq \Lambda$ , then  $x =_L y$  is a labeled condition

- if  $x_1, \dots, x_n \in \mathcal{X}$ ,  $P_L$  a labeled predicate, then  $P_L(x_1, \dots, x_n)$  is a labeled condition
- if  $\varphi$  and  $\psi$  are labeled conditions,  $L \subseteq \Lambda$ , then  $\neg_L \varphi$ ,  $\varphi \vee_L \psi$ , and  $\varphi \Rightarrow_L \psi$  are also labeled conditions
- if  $U$  is a set of labeled reference markers and  $Con$  a set of labeled conditions, then  $\langle U, Con \rangle$  is an LDRS
  - $\varphi =_{\text{def}} \langle U(\varphi), Con(\varphi) \rangle$
  - $U_L(\varphi) =_{\text{def}} \{x \mid \exists K [K \cap L \neq \emptyset \wedge x_K \in U(\varphi)]\}$

## Appendix II: semantics

- $\mathcal{M} = \langle D, W, R \rangle$ 
  - $D$  is a domain of individuals
  - $W$  is a set of extensional interpretation functions from basic predicates into subsets of  $D$
  - $R \subseteq W^2$
- $f[X]g =_{\text{def}} f \subseteq g \wedge Dom(g) = Dom(f) \cup X$
- $\|\varphi\|_{L,w}^f \downarrow =_{\text{def}} \|\varphi\|_{L,w}^f \in \{0, 1\}$
- $\|\varphi\|_{L,w}^f = \begin{cases} \{g \mid f[U_L(\varphi)]g \wedge \\ \wedge \forall \psi \in Con(\varphi) [\|\psi\|_{L,w}^g = 1]\} & \text{if } \exists g [f[U_L(\varphi)]g \wedge \\ \wedge \forall \psi \in Con(\varphi) [\|\psi\|_{L,w}^g \downarrow]] \\ \uparrow & \text{otherwise} \end{cases}$
- $\|x =_K y\|_{L,w}^f = \begin{cases} = 1 & \text{if } K \cap L = \emptyset \text{ or } x, y \in Dom(f) \wedge f(x) = f(y) \\ = 0 & \text{if } K \cap L \neq \emptyset \text{ and } x, y \in Dom(f) \wedge f(x) \neq f(y) \\ \text{undefined} & \text{otherwise} \end{cases}$
- $\|P_K(x_1, \dots, x_n)\|_{L,w}^f = \begin{cases} = 1 & \text{if } K \cap L = \emptyset \text{ or} \\ & x_1, \dots, x_n \in Dom(f) \wedge \langle f(x_1), \dots, f(x_n) \rangle \in w(P) \\ = 0 & \text{if } K \cap L \neq \emptyset \text{ and} \\ & x_1, \dots, x_n \in Dom(f) \wedge \langle f(x_1), \dots, f(x_n) \rangle \notin w(P) \\ \text{undefined} & \text{otherwise} \end{cases}$
- $\|\neg_K \psi\|_{L,w}^f = \begin{cases} = 1 & \text{if } K \cap L = \emptyset \text{ or } \|\psi\|_{K \cap L, w}^f = \emptyset \\ = 0 & \text{if } K \cap L \neq \emptyset \text{ and } \|\psi\|_{K \cap L, w}^f \downarrow \wedge \|\psi\|_{K \cap L, w}^f \neq \emptyset \\ \text{undefined} & \text{otherwise} \end{cases}$

- $\|\psi \vee_K \chi\|_{L,w}^f = \begin{cases} = 1 & \text{if } \|\psi\|_{K \cap L,w}^f \downarrow \wedge \|\chi\|_{K \cap L,w}^f \downarrow \wedge (\|\psi\|_{K \cap L,w}^f \cup \|\chi\|_{K \cap L,w}^f \neq \emptyset) \\ = 0 & \text{if } \|\psi\|_{K \cap L,w}^f = \|\chi\|_{K \cap L,w}^f = \emptyset \\ \text{undefined} & \text{otherwise} \end{cases}$
- $\|\psi \Rightarrow_K \chi\|_{L,w}^f = \begin{cases} = 1 & \text{if } \|\psi\|_{K \cap L,w}^f \downarrow \wedge \forall g \in \|\psi\|_{K \cap L,w}^f : \|\chi\|_{K \cap L,w}^g \downarrow \wedge \|\chi\|_{K \cap L,w}^g \neq \emptyset \\ = 0 & \text{if } \exists g \in \|\psi\|_{K \cap L,w}^f : \|\chi\|_{K \cap L,w}^g = \emptyset \\ \text{undefined} & \text{otherwise} \end{cases}$
- $\|\Box_K \psi\|_{L,w}^f = \begin{cases} = 1 & \text{if } \forall w' R w : \|\psi\|_{K \cap L,w'}^f \downarrow \wedge \|\psi\|_{K \cap L,w'}^f \neq \emptyset \\ = 0 & \text{if } \exists w' R w : \|\psi\|_{K \cap L,w'}^f = \emptyset \\ \text{undefined} & \text{otherwise} \end{cases}$

## References

- GEURTS, BART. 1997. Good news about the description theory of names. *Journal of Semantics* 14.319–348.
- . 1999. *Presuppositions and Pronouns*. Amsterdam: Elsevier.
- , & EMAR MAIER, ms. Layered DRT. (2003).
- KARTTUNEN, LAURI, & STANLEY PETERS. 1979. Conventional implicature. In *Syntax and Semantics II: Presupposition*, ed. by ChoonKyo Oh & David A. Dineen, 1–56. New York: Academic Press.
- MAIER, EMAR, & ROB A. VAN DER SANDT. to appear. Denial and correction in Layered DRT. In *Proc. of DiaBruck 2003*.
- VAN DER SANDT, ROB A. 1992. Presupposition projection as anaphora resolution. *Journal of Semantics* 9.333–377.